

EML: Linear Algebra over the Integers

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Preface

This article is the result of a three-months project about linear algebra over the integers. The programming part was completed using SageMath and its matrix related features.

This topic appeared particularly appealing as it seemed to allow us to find some nice results and tools while being only in first year. We thought it could bring us some interesting knowledge and useful skills that we would be able to reuse during the rest of our degree. We must say we were not disappointed at all so we will try to properly share our enjoyment while presenting our work in this article.

We wish to thank Prof. Dr. Gabor Wiese and Luca Notarnicola for their guidance throughout the making of this article.

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1 Introduction

Definition 1.1. Hermite Normal Form (HNF) We will say that a $m \times n$ matrix $M = (m_{i,j})$ with integer coefficients is in HNF if there exists $r \leq n$ and a strictly increasing map f from $[r+1, n]$ to $[1, m]$ satisfying the following properties :

- (1) For $r + 1 \leq j \leq n$, $m_{i,j} = 0$ if $i > f(j)$ and $0 \leq m_{f(k),j} < m_{f(k),k}$ if $k < j$.
- (2) The first r columns of M are equal to 0.

Remark. In the important special case where $m = n$, and $f(k) = k$ ($\iff \det(M) \neq 0$), M is in HNF if it satisfies the following conditions :

- (1) M is an upper triangular matrix, i.e. $m_{i,j} = 0$ if $i > j$.
- (2) For every i , we have $m_{i,i} > 0$.
- (3) For every $j > i$, we have $0 \leq m_{i,j} < m_{i,i}$.

Theorem 1.2. Let $A \in \mathcal{M}_{m,n}(\mathbb{Z})$. Then there exists a unique $m \times n$ matrix $B = (b_{i,j})$ in HNF of the form $B = AU$ with $U \in GL_n(\mathbb{Z})$ where $GL_n(\mathbb{Z})$ is the group of matrices with integer coefficients which are invertible, i.e. whose determinant is equal to ± 1 .

Proof. The proof is given by the algorithm in section 2.1 on page 3.

Remark. The reader should take note that some other literature may give a different definition for the HNF form of a matrix. This explains the possible difference between results obtained on HNF online calculators and the ones obtained in this paper for instance. For the entire article, a matrix will be in HNF if and only if it is in accordance with the definition 1.1.

Here are a few example to set down what being in HNF means. We will see a couple cases with different numbers of lines and columns:

Example 1.3.

a) Case of a 2×2 matrix : $M_1 = \begin{pmatrix} 5 & 7 \\ 2 & 3 \end{pmatrix}$

M_1 is not in HNF because it is a square matrix which is not upper triangular.

b) Case of a 2×5 matrix : $M_2 = \begin{pmatrix} 0 & 0 & 2 & 6 & 13 \\ 0 & 0 & 0 & 7 & 1 \end{pmatrix}$ is not in HNF because the first r columns from definition 1.1 are not equal to 0 (i.e. $M_{1,3} = 2 \neq 0$).

c) Case of a 4×4 matrix. : $M_3 = \begin{pmatrix} 4 & 8 & 5 & 11 \\ 7 & 2 & 6 & 1 \\ 1 & 17 & 7 & 12 \\ 9 & 3 & 5 & 8 \end{pmatrix}$ is not in HNF because it is a square matrix which is not upper triangular.

d) Case of a 3×5 matrix. : $M_4 = \begin{pmatrix} 0 & 0 & 51 & 50 & 11 \\ 0 & 0 & 0 & 28 & 2 \\ 0 & 0 & 0 & 0 & 13 \end{pmatrix}$ is in HNF because it fits the criteria from definition 1.1 (where $r = 2$).

2 A systematic HNF computation

In the first subsection, we will present the HNF algorithm in a spoken-like way. By doing so, we intend to making it easier for the reader to understand the procedure.

Then, in the following subsection we will present and explain the actual Sage-Math code and the functions used.

2.1 Algorithm

Algorithm 2.1. *Given an $m \times n$ A with integer coefficients $(a_{i,j})$, this algorithm finds the HNF W of A . Let us call $(w_{i,j})$ for the coefficients of W , A_i (resp. W_i) for the i^{th} column of A (resp. W)*

- Step 1.** [Initialize] Set $i \leftarrow m$, $k \leftarrow n$, $\begin{cases} l = 1 & \text{if } m \leq n \\ l = m - n + 1 & \text{if } m > n \end{cases}$
- Step 2.** [Row finished ?] If all the $a_{i,j}$ with $j < k$ are zero, then if $a_{i,k} < 0$, replace column A_k with $-A_k$ and go to step 5.
- Step 3.** [Choose non-zero entry] Pick among the non-zero $a_{i,j}$ for $j \leq k$ one with the smallest absolute value, say a_{i,j_0} . Then if $j_0 < k$, exchange column A_k with column A_{j_0} . In addition, if $a_{i,k} < 0$, replace column A_k by $-A_k$. Set $b \leftarrow a_{i,k}$.
- Step 4.** [Reduce] For $j = 1, \dots, k - 1$ do the following : set $q \leftarrow \lfloor a_{i,j}/b \rfloor$ and $A_j \leftarrow A_j - qA_k$. Then go to step 2.
- Step 5.** [Final reductions] Set $b \leftarrow a_{i,k}$. If $b = 0$, set $k \leftarrow k + 1$ and go to step 6. Otherwise for $j > k$ do the following: set $q \leftarrow \lfloor a_{i,j}/b \rfloor$, and $A_j \leftarrow A_j - qA_k$.
- Step 6.** [Finished?] If $i = l$, then for $j = 1, \dots, n - k + 1$ set $W_j \leftarrow A_{j+k-1}$ and terminate the algorithm. Otherwise, set $i \leftarrow i - 1$, $k \leftarrow k - i$ and go to step 2.

This algorithm terminates since the $|a_{i,k}|$ are strictly decreasing every time we return from **Step 2** to **Step 4**. It is clear that upon termination W will be in accordance to the HNF definition given in definition 1.1. Moreover, as we only used elementary columns operations on A , W is the HNF of A . \square

2.2 The code

The code which will be given in this subsection will be added as figures and the comments/explanations will be within the figures.

△ We would like to bring to the attention of the reader that Sagemath starts counting rows and columns from 0. It is crucial to keep this point in mind as it means a matrix of size $m \times n$, according to SageMath, has rows numbered from 0 to $m-1$ and columns numbered from 0 to $n-1$.

One symbol “#” will be added before the first line of a comment (this is the normal way to add a comment on SageMath). But if a comment goes through several lines, we will use the following convention: two “#” for the second line, three “#” for the third line and so on. Here is an example to clarify the convention :

```
#comment number 1 (1st line)
#comment number 2 (1st line)
# #comment number 2 (2nd line)
#comment number 3 (1st line)
```

```

1 #Let A be a mxn matrix
2 ##this is a "random" matrix to set ideas down.
3
4 A=matrix([[21, 2, 3, 4, 5, 6],[7, 8, 9, 10, 11, 12],[13, 14, 15, 16, 17, 18]])
5 print A
6 i=A.nrows()-1 #set i as last row of matrix
7 k=A.ncols()-1 #set k as last column of matrix
8
9 #Initialise l:
10 if A.nrows()<=A.ncols():
11     l=0
12 else:
13     l=i-k
14 j=0
15
16 #Here starts the HNF Process:
17 #Step 1
18 while i!=1:
19
20     #We define a row-matrix r such that it has all the elements A[i][j] of the i-th row, with j<k
21     r=A.matrix_from_rows_and_columns([i],range(0,k))
22     print "This is r"
23     print r
24
25     #Step 2
26     while r.is_zero()!=True:
27         #Define a row-matrix y that has the absolute value of all the
28         #elements A[i][e] of the i-th row, with: e<=k & A[i][e] !=0
29         e=0
30         y=[]
31         while e<=k:
32             if A[i][e]!=0:
33                 y.append(abs(A[i][e]))
34             e=e+1
35         print "This is y"
36         print y
37
38         #Step 3
39         #Find minimum:
40         z=0
41         if len(y)>=2:
42             jo=min(y[z],y[z+1]) #if len(y)=2 we take the minimum of the two elements of y
43                                 #We still take the minimum of the first two elements if len(y)>2 (see next while loop)
44                                 #We call the minimum jo
45             z=2
46             while z<=len(y)-1: #This is a while loop that will find the minimum of y for the case len(y)>2
47                 jo=min(jo,y[z]) #as we found the minimum of the first two (called jo)
48
49                                 #we overwrite jo by saying that is the minimum of (the minimum of the elements
50                                 #before) jo and the z-th element
51             z=z+1
52         else:
53             if len(y)==1: #If y has only one element the minimum is the element itself
54                 jo=y[0]
55             print "this jo "
56             print jo

```

Figure 1: The entire HNF code (part 1/4).


```

57
58     #Now we search the column u to which jo belongs (from the 0-th to the (n-1)-th)
59     u=0
60     while abs(A[i][u])!=jo:
61         u=u+1
62     #if u is different from k we exchange column u and column k.
63     if abs(A[i][k])!= jo:
64         if u<k:
65             A.swap_columns(u,k)
66
67     #column Ak=-Ak if aik is negative
68     if A[i][k] < 0:
69         A.rescale_col(k,-1)
70     print "new A"
71     print A
72
73     #define b=a_{ik} (usually from all the steps before: b is smallest
74     ##positive integer, but b can also be 0 in some cases )
75     b = A[i][k]
76
77
78     #Step 4
79     z=0
80     #Az= Az -q*Ak for all integer 0<=z<k
81     while z <= k-1 :
82         q = round(A[i][z]/b)
83         print "this is q1"
84         print q
85         A.set_column(z,A.column(z)-q*A.column(k))
86         z = z+1
87     print A
88
89     #As this is a loop we need to redefine r such that we can check again the condition
90     r=A.matrix_from_rows_and_columns([i],range(0,k))
91     print "This is r after checking the condition (see comments)"
92     print r
93
94     #If r==0 we change column Ak=-Ak if a_{ik} is negative.
95     if (r.is_zero()==True)or(k==0) :
96         if A[i][k] < 0:
97             A.rescale_col(k,-1)
98         b = A[i][k]
99
100
101     #Step 5:
102     if b==0:
103         k=k+1
104
105     else:
106         #check if there exists at least one column after column k
107         if k < A.ncols()-1 :
108             w=k+1
109             # Aj<- Aj-q*Ak for j>k (not the same q as before )
110             while w<=A.ncols()-1:
111                 q = floor(A[i][w]/b)
112                 A.add_multiple_of_column(w,k,-q)
113                 w=w+1

```

Figure 2: The entire HNF code (part 2/4).

```

114     k=k-1 #as the whole algorithm is a while-Loop we need to redo this step until l=i
115     i=i-1]
116 #This happens when i=l, as i=l is the condition to break the Loop
117 ##above, but we still need to check the Last step one Last Time
118
119 #Final Stage
120 if i==l:
121     print "Final Stage"
122     #We define a row-matrix r such that its has all the elements A[i][j] of the i-th rows, with j<k
123     r=A.matrix_from_rows_and_columns([i],range(0,k))
124     print"This is r"
125     print r
126
127 #Step 2
128 while r.is_zero()!=True:
129     #Define a list/row-matrix y that has the absolute value of all the elements
130     ##A[i][e] of the i-th row, with: e<=k & A[i][e] !=0
131     e=0
132     y=[]
133     while e<=k:
134         if A[i][e]!=0:
135             y.append(abs(A[i][e]))
136         e=e+1
137     print "This is y"
138     print y
139
140 #Find minimum(Step3):
141 z=0
142 if len(y)>=2:
143     jo=min(y[z],y[z+1]) #if len(y)=2 we take the minimum of the two elements of y
144                        #We still take the minimum of the first two elements if len(y)>2
145                        #We call the minimum jo
146     z=2
147     while z<=len(y)-1: #Here is a loop that will find the minimum of y for the case len(y)>2
148         jo=min(jo,y[z]) #as we found the minimum of the first two (called jo)
149                        #we overwrite jo by saying that is the minimum of
150                        ##(the minimum of the elements before) jo and the z-th element
151         z=z+1
152 else:
153     if len(y)==1: #If y has only one element the minimum is the element itself
154         jo=y[0]
155     print "this is jo "
156     print jo
157 #Next step we search the column u which has that element jo (from left to right)
158 u=0
159 while abs(A[i][u])!=jo:
160     u=u+1
161 #if u is different from k we exchange column u and column k
162 if abs(A[i][k])!= jo:
163     if u<k:
164         A.swap_columns(u,k)
165
166 #column Ak=-Ak if a_{ik} is negative.
167 if A[i][k] < 0:
168     A.rescale_col(k,-1)
169     print "new A"
170     print A

```

Figure 3: The entire HNF code (part 3/4).

```

172     #define b=aik (usually from all the steps before: b is smallest
173     ##positive integer, but b can also be 0 in some cases )
174     b = A[i][k]
175     #Step 4
176     z=0
177     #Az= Az -q*Ak for all integer 0<=z<k
178     while z <= k-1 :
179         q = round(A[i][z]/b)
180         print "this is q1"
181         print q
182         A.set_column(z,A.column(z)-q*A.column(k))
183         z = z+1
184     print A
185
186     #As this is a loop we need to redefine r st: we can check again the condition
187     r=A.matrix_from_rows_and_columns([i],range(0,k))
188     print "This is r after some computations"
189     print r
190
191     #If r==0 we change column Ak=-Ak if aik is smaller than 0
192     if (r.is_zero()==True)or(k==0) :
193         if A[i][k] < 0:
194             A.rescale_col(k,-1)
195         b = A[i][k] #b=aik
196         #Step 5:
197         if b==0:
198             k=k+1
199
200         else:
201             #check if there exists atleast one column after column k
202             if k < A.ncols()-1 :
203                 w=k+1
204                 # Aj<- Aj-q*Ak for j>k (not the same q as before )
205                 while w<A.ncols()-1: #Colonne pour ligne
206                     q = floor(A[i][w]/b)
207                     A.add_multiple_of_column(w,k,-q)
208                     w=w+1
209     #Define W the hermit-normal-form of A that starts from the k-th column as everything before is 0
210     W=A.matrix_from_columns(range(k,A.ncols()))
211     print "This is W"
212     print W

```

Figure 4: The entire HNF code (part 4/4).

2.3 Applications using the HNF

The process computing the HNF of a matrix can be very useful in several domains. In this subsection we will give some of these domains with a brief explanation of the use in each case.

2.3.1 Linear Diophantine equations

The HNF appears to be very useful for solving systems of linear Diophantine equations.

Let A be a matrix and b a vector. We are looking for a vector x such that $Ax = b$. In fact, we may assume that A has full row rank; otherwise, we may remove redundant equations from the system. Yet, we assume all input

data to be rational.

We can find a unimodular matrix U such that $[H \mid 0] = AU$ is a matrix in Hermite normal form. Now, we can apply a standard trick to transform a system of linear equations (and actually, a system of linear inequalities, too) into a more suitable form: $Ax = b$ is equivalent to $(AU)(U^{-1}x) = b$, and therefore, $[H \mid 0]z = b$, where $z = U^{-1}x$ is integral if and only if x is integral. We can observe that the last components of z may take arbitrary values, while feasibility of the system $[H \mid 0]z = b$, and therefore, of $Ax = b$, depends only on whether the vector $H^{-1}b$ is integral.

From there, we can find ways to set conditions for a system of linear Diophantine equations to have a solution.

We will not go any further into explaining how the HNF is used to solve Linear Diophantine equations as this subsection is just meant to give leads on some applications. Further information can be found on some of the references listed at the end of this article.

2.3.2 case $Ax=0$

The HNF can also be used to determine the kernel of a matrix. Here is the code to do so :

```

1 A=matrix([[21, 2, 3],[7, 8, 9],[13, 14, 15]])#Define a Matrix mxn
2 print A
3 #Initialise:
4 i=A.nrows()-1 #i=m
5 j=A.ncols()-1 #j(=r)=n
6 k=A.ncols()-1
7 U=matrix.identity(A.ncols()) #U=In(Identity matrix nxn)
8 #set l:
9 if A.nrows()<=A.ncols():
10     l=0
11 else:
12     l=i-k
13 #Ax=0 Process:
14 while i!=l:
15     #We finish when i=l (l depending on matrix A)
16     while j!=0:
17         #Algorithm until j=0
18         j=j-1
19         if A[i][j]!=0:
20             #We compute xgcd= extended Euclid Algorithm
21             a=xgcd(A[i][k],A[i][j])
22             d=a[0] #d=gcd(aij,aik)
23             u=a[1] #u first integer
24             v=a[2] #v second integer
25             #Computations on columns of A and U
26             #We want to exchange values(The reason we create the temporary vector B)
27             B=vector(u*A.column(k)+v*A.column(j)) #B defined such that we have for i-th row the gcd of column j and k of A
28             #We want to exchange Values (The reason we create the temporary vector C)
29
30             #Important Step: Inorder to keep the values of aik and aij we first compute this step on U and then on A
31
32             C=vector(u*U.column(k)+v*U.column(j)) #C defined as B in the Beginning but for our matrix U
33             U.set_column(j,(A[i][k]/d)*U.column(j)-(A[i][j]/d)*U.column(k)) #Uj=aik/d*Uj-aik/d*Uk
34             U.set_column(k,C) #Uk=C
35             A.set_column(j,(A[i][k]/d)*A.column(j)-(A[i][j]/d)*A.column(k)) #set Aj <- aik/d *Aj - aij/d*Ak
36             A.set_column(k,B) #Ak = B
37             print "New A"
38             print A
39             print "U:"
40             print U
41
42
43 #if aik<0 Ak=-Ak, Uk=-Uk
44 if A[i][k] < 0:
45     A.rescale_col(k,-1)
46     U.rescale_col(k,-1)
47 #We define b = aik
48 b=A[i][k]
49 #check if b=0 or not
50 if b==0:
51     k=k+1
52 else:
53     #check if there exists atleast one column after column k
54     if k<A.ncols()-1 :
55         r=k+1

```

Figure 5: The kernel code (part 1/2).

```

56     # Aj<- Aj-q*Ak for j>k (here j is r)
57     while r<-A.ncols()-1: #Colonne pour ligne
58         q = floor(A[i][r]/b)
59         A.add_multiple_of_column(r,k,-q)
60         U.add_multiple_of_column(r,k,-q)
61         r=r+1
62     #As this a while loop we define it like this
63     i=i-1
64     k=k-1
65     j=k
66 #Final Step: i==l
67
68 if i==l:
69     while j!=0: #Algorithm until j=0
70         if A[i][j]==0: #if aij =0 we "move on"
71             j=j-1
72         else: #otherwise we also set j=j-1 like before
73             j=j-1
74             #We compute xgcd= extended Euclid Algorithm
75             a=xgcd(A[i][k],A[i][j])
76             d=a[0] #d=gcd(aik,aij)
77             u=a[1] #u first integer
78             v=a[2] #v second integer
79             #Computations on columns of A and U
80             #We want to exchange values(The reason we create the temporary vector B)
81             B=vector(u*A.column(k)+v*A.column(j)) #B defined such that we have for i-th row d (gcd(aik,aij))
82             #We want to exchange Values (The reason we create the temporary vector C)
83             C=vector(u*U.column(k)+v*U.column(j)) #C defined as B in the Beginning but for our matrix U
84             U.set_column(j,(A[i][k]/d)*U.column(j)-(A[i][j]/d)*U.column(k)) #Uj=aik/d*Uj-aij/d*Uk
85             U.set_column(k,C) #Uk=C
86             A.set_column(j,(A[i][k]/d)*A.column(j)-(A[i][j]/d)*A.column(k)) #set Aj <- aik/d *Aj - aij/d*Ak
87             A.set_column(k,B) #Ak = B
88             print A
89             print "U:"
90             print U
91     #if aik<0 Ak=-Ak, Uk=-Uk
92     if A[i][k] < 0:
93         A.rescale_col(k,-1)
94         U.rescale_col(k,-1)
95     #We define b = aik
96     b=A[i][k]
97     #check if b=0 or not
98     if b==0:
99         k=k+1
100     else:
101         #check if there exists atleast one column after column k
102         if k<A.ncols()-1 :
103             r=k+1
104             # Aj<- Aj-q*Ak for j>k (here j is r)
105             while r<-A.ncols()-1: #Colonne pour ligne
106                 q = floor(A[i][r]/b)
107                 A.add_multiple_of_column(r,k,-q)
108                 U.add_multiple_of_column(r,k,-q)
109                 r=r+1
110     #We define a basis of the kernel of A with kerA=Uj and j=0,...,k (1,...,k-1)
111     kerA=U.matrix_from_columns(range(0,k))
112     print "This is the kernel of A"
113     print kerA
114

```

Figure 6: The kernel code (part 2/2).

References

- https://en.wikipedia.org/wiki/Hermite_normal_form
- <https://disopt.epfl.ch/webdav/site/disopt/shared/IntPoints2009/lattices.pdf>
- <https://disopt.epfl.ch/webdav/site/disopt/shared/IntPoints2009/hnf.pdf>
- https://ac.els-cdn.com/0024379590902285/1-s2.0-0024379590902285-main.pdf?_tid=b6c4503c-cd0e-45c6-8050-f884bfe51abc&acdnat=1528063809_91dde5a4a0594a39b55b29edfce05eb7
- <https://www.uow.edu.au/thomaspl/pdf/Tourloupis13.pdf>

Appendix

This appendix is meant to explain the SageMath built-in functions used in the HNF code (section 2.2) and/or in the code used to find the Kernel of a matrix.

Let A be a matrix of size $(m \times n)$.

A.nrows() : gives the number of rows the matrix A has
(= m)

A.ncols() : gives the number of columns the matrix A
has (= n)

A.matrix_from_rows_and_columns($[i_1, i_2, \dots, i_p], [j_1, j_2, \dots, j_q]$)
: creates a submatrix of A created from the rows and columns selected. Where $0 \leq i_a \leq (m - 1)$ is the a -th line and $0 \leq j_b \leq (n - 1)$ is the b -th column. Lines and Columns can be reordered or repeated.

A.matrix_from_columns($[j_1, j_2, \dots, j_q]$) : command `A.matrix_from_rows_and_columns` (see before) where all rows are kept and stay in the same order.

A.swap_columns(u,k) : exchanges the column u and column k with $0 \leq u, k \leq m - 1$.

A.rescale_col(k,s) : multiplies the column k ($0 \leq k \leq m - 1$) by s . This command only works if s matches

the ring of A .

Example : ring of $A = \mathbb{ZZ}$ and $s = 0.5 \Rightarrow ERROR$

A.set_column(i,u) : set the i -th column to u with u being a vector of size $(n \times 1)$.

xgcd(i,j) : Computes the Extended Euclidean Algorithm to find coefficients u and v such that: $u \times i + v \times j = gcd(i, j)$ returns $(gcd(i, j), u, v)$. Caution: $|u|$ and $|v|$ are not minimal.

A.add_multiple_of_column(r,k,q) : replaces the r -th column by $r + q \times k$ where q belongs to A 's ring and k stands for the k -th column ($0 \leq k \leq m - 1$).

A.is_zero() : returns true if matrix A is the zero matrix.