# Modular Forms Project: Computing Expansions at Cusps

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## 1 Answers and Programs

#### Exercise 1.

(1). Let as usual  $T = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ . If  $\gamma = \begin{pmatrix} A & B \\ C & D \end{pmatrix} \in \Gamma$ , we compute that

$$\gamma T^m \gamma^{-1} = \begin{pmatrix} 1 - mAC & mA^2 \\ -mC^2 & 1 + mAC \end{pmatrix} .$$

This belongs to  $\Gamma_0(N)$  if and only if  $N \mid mC^2$ , hence if and only if  $w_0 \mid m$  with  $w_0 = N/\gcd(C^2, N)$ . Since  $F \in M_k(\Gamma_0(N), \chi)$ , we thus have  $F \mid_k \gamma T^{w_0 n} \gamma^{-1} = \chi(1 + ACw_0 n)F$ , or

$$(F|_k\gamma)(\tau + w_0n) = \chi(1 + ACw_0n)(F|_k\gamma)(\tau) .$$

If we choose  $n = H = w_1/w_0$ , we note that

$$1 + ACw_0 n = 1 + ACw_1 = 1 + A(C/\gcd(C, N))N \equiv 1 \pmod{N},$$

so  $\chi(1 + ACw_0 n) = 1$  and  $(F|_k \gamma)(\tau + w_1) = (F|_k \gamma)(\tau)$ .

(2). It follows from the trivial equality  $(\gamma T^{w_0} \gamma^{-1})^n = (\gamma T^{w_0} \gamma^{-1})$  that for all *n* we have  $\chi(1 + ACw_0 n) = \chi(1 + ACw_0)^n$ . Applying this to  $n = H = w_1/w_0$ shows that  $\chi(1 + ACw_0)$  is an *H*th root of unity, hence uniquely of the form  $e^{2\pi i \alpha/H}$  for some integer  $\alpha$  such that  $0 \leq \alpha < H$ . It is clear that if H = 1 we have  $\alpha = 0$ , but also if  $\chi$  is a trivial character.

(3). If we change  $\tau$  into  $\tau + w_0$  in the expansion of  $F|_k \gamma$  we have on the one hand

$$(F|_k\gamma)(\tau+w_0) = \sum_{n\geq 0} a_\gamma(n) e^{2\pi i n/H} q^{n/w_1},$$

and on the other hand

$$(F|_k\gamma)(\tau+w_0) = e^{2\pi i\alpha/H}F|_k\gamma(\tau) .$$

Identifying coefficients of  $q^{n/w_1}$  gives

$$a_{\gamma}(n)(e^{2\pi i n/H} - e^{2\pi i \alpha/H}) = 0$$
,

which is equivalent to  $a_{\gamma}(n) = 0$  when  $n \not\equiv \alpha \pmod{H}$ . It follows that we can write

$$(F|_{\gamma})(\tau) = q^{\alpha/w_1} \sum_{n \ge 0} b(n)q^{n/w_0}$$

with  $b(n) = a_{\gamma}(nH + \alpha)$ .

### Exercise 2.

(1). Multiplying by a common denominator M of A, B, C, and D, we can assume that  $\gamma \in M_2^+(\mathbb{Z})$  (we will simply divide the triangular matrix by M at the end). We want to find  $\delta = \begin{pmatrix} U & V \\ W & X \end{pmatrix} \in \Gamma$  and a, b, and d such that  $\gamma = \delta \begin{pmatrix} a & b \\ 0 & d \end{pmatrix}$ . This implies first that Ua = A and Wa = C, and since U and W must be coprime, up to sign this has the only solution  $a = g := \gcd(A, C)$ , U = A/g, W = C/g.

Thus U and W are indeed coprime, hence by Euclid there exist (nonunique) V and X such that UX - VW = 1. The other condition to be verified is  $\delta(b,d)^t = (B,D)^t$ , and since  $\delta \in \Gamma$  this gives

$$\begin{pmatrix} b \\ d \end{pmatrix} = \begin{pmatrix} X & -V \\ -W & U \end{pmatrix} \begin{pmatrix} B \\ D \end{pmatrix} ,$$

in other words b = BX - DV, d = DU - BW = (AD - BC)/g.

This last equality shows that d does not depend on the choice of V and X (which is trivial since the determinant of the triangular matrix must be equal to AD - BC), so only b does. If (V, X) is changed into  $(V + \lambda U, X + \lambda W)$  then b is changed into  $b + \lambda(BW - DU) = b - \lambda d$ , hence  $\begin{pmatrix} a & b \\ 0 & d \end{pmatrix}$  into  $T^{-\lambda} \begin{pmatrix} a & b \\ 0 & d \end{pmatrix}$  and  $\delta$  into  $\delta T^{\lambda}$ .

As mentioned above, we finish by dividing a, b, and d by M.

(2) The matrix  $\begin{pmatrix} a & b \\ 0 & d \end{pmatrix}$  corresponds to the Möbius transformation  $(a\tau + b)/d = (a/d)\tau + (b/d)$ . Thus, if  $G := F|_k \delta = \sum_{n>0} a_\delta(n)q^{n/w_1}$  we have

$$G(\tau + b/d) = \sum_{n \ge 0} a_{\delta}(n) e^{2\pi i (b/d)n/w_1} q^{n/w_1} ,$$

hence

$$G((a/d)\tau + b/d) = \sum_{n \ge 0} a_{\delta}(n) e^{2\pi i(b/d)n/w_1} q^{an/(dw_1)} .$$

#### Exercise 3.

(1). Note that Pari/GP already has a built-in command mfslashexpansion to compute  $F|_k\gamma$ . Of course we want to implement Collins's algorithm, so we do not want to use Pari/GP's command. However, handling general Dirichlet characters, although not difficult is slightly tedious, so we will simply use mfslashexpansion to compute the integer  $\alpha$  such that  $0 \leq \alpha < H$  which is needed in the expansions, which is not really cheating. We could of course also restrict the algorithm to the known cases where  $\alpha = 0$ , i.e.,  $\chi$  trivial or  $H = \gcd(C^2, N)/\gcd(C, N) = 1$ .

A possible program is as follows:

```
/* Collins's algorithm for cusp expansions. */
collins(mf,F,ga,K0,G0=1)=
```

```
ſ
  my(N,k,A,B,C,D,E,K,G,L,w0,w1,H,al,T,im);
  my(Y,S,Q,R,M,MS,X);
  [N,k]=mfparams(mf); [A,C]=ga[,1]; [B,D]=ga[,2];
  E=bitprecision(1.)*0.69315;
  if (K0*GO<E,G=GO;K=ceil(E/G),K=KO;G=E/K);</pre>
  L=2*K; wO=N/gcd(C*C,N); w1=N/gcd(C,N); H=w1/w0;
  mfslashexpansion(mf,F,ga,0,0,&params); al=params[1]*w1;
/* only used to compute alpha. */
  T=vector(L); /* vector of random tau's. */
  im=G*w0*I/(2*Pi); /* desired imaginary part */
  for (i=1,L,T[i]=(random(1.)-1/2)*w0+im-D/C);
/* random real parts in the interval [-D/C-w0/2,-D/C+w0/2]. */
  Y=vectorv(L,i,(C*T[i]+D)^(-k)*mfeval(mf,F,(A*T[i]+B)/(C*T[i]+D)));
  S=vector(L,i,exp(2*Pi*I*T[i]/w1));
  Q=vector(L,i,S[i]^H); R=vector(L,i,S[i]^al);
  M=matrix(L,K,i,j,R[i]*Q[i]^(j-1));
  MS=conj(M~); X=matsolve(MS*M,MS*Y);
  X[1..K0];
}
```

In the above program there is an additional cheat: we use the Pari/GP function mfeval instead of computing sufficiently many terms of the Fourier expansion of F at infinity and summing the corresponding series.

(2). For  $E_4$  we can simply write

```
E4=mfEk(4); mf=mfinit(E4);
C=collins(mf,E4,[0,-1;1,0],20)
CI=bestappr(C,10<sup>4</sup>)
CI<sup>~</sup>==mfcoefs(E4,19)
norml2(C-CI)
```

We first obtain some complex numbers which are clearly close to rational numbers (in fact to integers here). We then compute the best rational approximations to these complex numbers with denominator less than  $10^4$ , and we recognize the coefficients of  $E_4$ , which we can check without reading all coefficients by the command  $CI^{-=mfcoefs}(E4, 19)$  (note that CI is a column vector, so that it must be transposed, and that mfcoefs(F,L) gives coefficients a(0) up to a(L), so to obtain 20 coefficients we must choose L = 19).

Finally, to check how close the least squares approximation was, we compute the square of the  $L^2$  norm of the difference of C and CI (note that contrary to what its name could imply, the **norm12** command gives the *square* of the  $L^2$  norm). We obtain  $1.076 \cdots E - 45$ , so the values of the coefficients have approximately 22 correct decimal digits, the computation being done with the default accuracy of 38.

A similar program (replacing the first line by D=mfDelta(); mf=mfinit(D,1); ) also gives correct results, but since the weight is much larger, we only obtain 20 correct decimal digits.

(3). (a). To initialize the computation, as required we first write:

```
mf=mfinit([6,4],0); F=mfeigenbasis(mf)[1];
ga=[1,-1;3,-2]; CF=mfcoefs(F,19)
C=collins(mf,F,ga,20);
CI=bestappr(C,10<sup>4</sup>)
CI<sup>~</sup>==CF/4
norml2(C-CI)
```

which first outputs

[0, 1, -2, -3, 4, 6, 6, -16, -8, 9, -12, 12, -12, 38, 32, -18, 16, -126, -18, 20]

as the first 20 Fourier coefficients of F, then the complex numbers computed by the program are approximated by

[0, 1/4, -1/2, -3/4, 1, 3/2, 3/2, -4, -2, 9/4, -3, 3, -3, 19/2, 8, -9/2, 4, -63/2, -9/2, 5]

The next command checks that this is indeed 1/4 times the coefficients of F, and the last command (which outputs 7.088E - 58) shows that the algorithm gives values which are correct to 29 decimal digits.

(3). (b). To check the eta product without too much using Pari/GP's functionalities, the simplest is to do as follows:

```
G=mffrometaquo([1,2;2,2;3,2;6,2]); par=mfparams(G)[1..2]
st=mfsturm(par)
mfcoefs(F,st)==mfcoefs(G,st)
```

The first line outputs [6, 4], which tells us that  $G \in M_4(\Gamma_0(6))$ , the second line gives the *sturm bound*, here 5, and the last command checks that the first st+1 (counting the constant term) coefficients of F and G are equal, thus *proving* that F = G. The shorter Pari/GP command mfisequal(F,G) does exactly that.

(3). (c). The builtin mfslashexpansion(mf,F,ga,19,1) command outputs [0, 1/4, -1/2, -3/4, 1, 3/2, 3/2, -4, -2, 9/4, -3, 3, -3, 19/2, 8, -9/2, 4, -63/2, -9/2, 5]

equal to what we found above.

(4). We write the following simple programs:

1/2: 4.732342396833962897 E-64 (time = 6757)
1: 2.763570679490571905 E-57 (time = 945)
2: 2.1521231746487435604 E-44 (time = 232)
3: 2.1359265017258375511475323561759823231 E-28 (time = 132)
4: 1.2676095444665378389804436912111348768 E-12 (time = 64)

This output shows the (logical) fact that as  $G_0$  increases the time becomes considerably shorter, but the accuracy considerably worse (only 6 correct decimal digits for  $G_0 = 4$ ).

(5). (a). As required, we write

```
mf=mfinit([27,4],0); F=mfeigenbasis(mf)[2];
ga=[1,-1;3,-2]; CF=mfcoefs(F,19)
C=collins(mf,F,ga,20);
```

The Fourier coefficients of F are thus

[0, 1, -3, 0, 1, -15, 0, -25, 21, 0, 45, 15, 0, 20, 75, 0, -71, -72, 0, 2]

We do not print the output of the collins command since they are complicated complex numbers. Instead, as suggested we write the following programs which generalises the **bestappr** function to cyclotomic fields by using **lindep** instead:

```
/* Find z in $\Q(e^{2\pi i/N})$. */
cycloapprox(z,N)=
{
    my(pn=eulerphi(N)-1,P=powers(exp(2*I*Pi/N),pn),T);
    T=powers(t,pn);
    L=lindep(concat(-z,P));
    return (T*(L[2..#L]/L[1]));
}
```

cyclovecapprox(V,N)=apply(z->cycloapprox(z,N),V);

If we execute cyclovecapprox(C,9), we obtain

```
[0, -1/9*t<sup>4</sup>, -1/3*t<sup>5</sup> - 1/3*t<sup>2</sup>, 0, 1/9*t<sup>4</sup> + 1/9*t, 5/3*t<sup>2</sup>,
0, 25/9*t, -7/3*t<sup>5</sup>, 0, -5*t<sup>4</sup>, 5/3*t<sup>5</sup> + 5/3*t<sup>2</sup>, 0,
20/9*t<sup>4</sup> + 20/9*t, -25/3*t<sup>2</sup>, 0, 71/9*t, 8*t<sup>5</sup>, 0, -2/9*t<sup>4</sup>]
```

(output edited for clarity), where in the above t is to be interpreted as  $e^{2\pi i/9}$ . In fact, since  $t^3 + 1 = -t^6$ , it is easy to see that all the coefficients are rational multiples of 9th roots of unity.

(5). (b). Here we must write a generalization of the program given in the previous question. For this, we simply replace the tstg program by:

```
tstg(G0)=
{
    my(C=collins(mf,F,ga,20,G0));
    CI=cyclovecapprox(C,9);
    norml2(C-subst(CI,t,exp(2*Pi*I/9)));
}
```

and the rest is the same. The output will now be

```
1/2: 8.111206046459227532 E-64 (time = 10292)
1: 1.7354873000414302779 E-57 (time = 2689)
2: 3.3933846210407921667 E-46 (time = 828)
3: 1.6704210168908523895065512908299597748 E-46 (time = 585)
4: 1.4631342795842769204704557109803552715 E-45 (time = 617)
```

Here there is much less loss of accuracy when  $G_0$  increases.

(5). (c). Typing CJ=mfslashexpansion(mf,F,ga,19,1) of course gives the same expansion, except that the coefficients are expressed as POLMODS in the field defined by  $t^6 + t^3 + 1 = 0$ , i.e., the ninth cyclotomic field, so to obtain exactly the same output one simply writes lift(CJ).