
Computing congruences of modular forms modulo prime powers

Gabor Wiese

(joint work with Xavier Taixés i Ventosa)

Institut für Experimentelle Mathematik

Universität Duisburg-Essen

16 November 2009

Plan

- (I) Congruences mod ℓ^n .
- (II) Computing them.
- (III) Modular forms mod ℓ^n .
- (IV) Examples and applications.

Congruences mod ℓ^n (black board).

Next: Computing congruences mod ℓ^n .

Computing congruences mod ℓ^n

Problem: Let $P, Q \in \mathbb{Z}[X]$ be monic coprime polynomials.

For which prime powers ℓ^n are there $\alpha, \beta \in \overline{\mathbb{Z}}$ such that

- (i) $P(\alpha) = Q(\beta) = 0$ and
- (ii) $\alpha \equiv \beta \pmod{\ell^n}$?

Partial solution (very fast):

Reduced resultant = congruence number (global).

Complete solution:

Newton polygon method (local).

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

Sylvester map:

$$\begin{array}{ccc} \mathbb{Z}[X]_{} & \times & \mathbb{Z}[X]_{} \\ \{X^{v-1}, \dots, X, 1\} & & \{X^{u-1}, \dots, X, 1\} \end{array} \xrightarrow{(r,s) \mapsto rP + sQ} \mathbb{Z}[X]_{} \quad \{X^{u+v-1}, \dots, X, 1\}.$$

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

Sylvester map:

$$\begin{array}{ccc} \mathbb{Z}[X]_{} & \times & \mathbb{Z}[X]_{} \\ \{X^{v-1}, \dots, X, 1\} & & \{X^{u-1}, \dots, X, 1\} \end{array} \xrightarrow{(r,s) \mapsto rP + sQ} \mathbb{Z}[X]_{} \quad \{X^{u+v-1}, \dots, X, 1\}.$$

Sylvester matrix (for column vectors) with $u = 3$ and $v = 2$:

$$\begin{pmatrix} a_0 & 0 & b_0 & 0 & 0 \\ a_1 & a_0 & b_1 & b_0 & 0 \\ a_2 & a_1 & b_2 & b_1 & b_0 \\ a_3 & a_2 & 0 & b_2 & b_1 \\ 0 & a_3 & 0 & 0 & b_2 \end{pmatrix}$$

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

$$\begin{pmatrix} a_0 & 0 & b_0 & 0 & 0 \\ a_1 & a_0 & b_1 & b_0 & 0 \\ a_2 & a_1 & b_2 & b_1 & b_0 \\ a_3 & a_2 & 0 & b_2 & b_1 \\ 0 & a_3 & 0 & 0 & b_2 \end{pmatrix}$$

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

$$\begin{pmatrix} a_0 & 0 & b_0 & 0 & 0 \\ a_1 & a_0 & b_1 & b_0 & 0 \\ a_2 & a_1 & b_2 & b_1 & b_0 \\ a_3 & a_2 & 0 & b_2 & b_1 \\ 0 & a_3 & 0 & 0 & b_2 \end{pmatrix}$$

Want to know its image for the basis $\{X^{u+v-1}, \dots, X, 1\}$.

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

$$\begin{pmatrix} a_0 & 0 & b_0 & 0 & 0 \\ a_1 & a_0 & b_1 & b_0 & 0 \\ a_2 & a_1 & b_2 & b_1 & b_0 \\ a_3 & a_2 & 0 & b_2 & b_1 \\ 0 & a_3 & 0 & 0 & b_2 \end{pmatrix}$$

Want to know its image for the basis $\{X^{u+v-1}, \dots, X, 1\}$.

May multiply by invertible integer matrices *from the right*.

I.e. may perform integral column operations.

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

$$\begin{pmatrix} a_0 & 0 & b_0 & 0 & 0 \\ a_1 & a_0 & b_1 & b_0 & 0 \\ a_2 & a_1 & b_2 & b_1 & b_0 \\ a_3 & a_2 & 0 & b_2 & b_1 \\ 0 & a_3 & 0 & 0 & b_2 \end{pmatrix} \circ \begin{pmatrix} * & * & * & * & * \\ * & * & * & * & * \\ * & * & * & * & * \\ * & * & * & * & * \\ * & * & * & * & * \end{pmatrix} =$$

Want to know its image for the basis $\{X^{u+v-1}, \dots, X, 1\}$.

May multiply by invertible integer matrices *from the right*.

I.e. may perform integral column operations.

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

$$= \begin{pmatrix} * & 0 & 0 & 0 & 0 \\ * & * & 0 & 0 & 0 \\ * & * & * & 0 & 0 \\ * & * & * & * & 0 \\ * & * & * & * & c \end{pmatrix}$$

Want to know its image for the basis $\{X^{u+v-1}, \dots, X, 1\}$.

Congruence number

$$P(X) = \sum_{k=0}^u a_k X^{u-k}, \quad Q(X) = \sum_{k=0}^v b_k X^{v-k} \in \mathbb{Z}[X].$$

$$= \begin{pmatrix} * & 0 & 0 & 0 & 0 \\ * & * & 0 & 0 & 0 \\ * & * & * & 0 & 0 \\ * & * & * & * & 0 \\ * & * & * & * & c \end{pmatrix}$$

Want to know its image for the basis $\{X^{u+v-1}, \dots, X, 1\}$.

Congruence number $c(P, Q)$ is the bottom right entry!

It divides the resultant of P, Q (determinant of $S(P, Q)$).

Congruence number

$$P(X) = X - a, \quad Q(X) = X - b.$$

$$S(P, Q) = \begin{pmatrix} 1 & 1 \\ -a & -b \end{pmatrix}$$

Congruence number

$$P(X) = X - a, \quad Q(X) = X - b.$$

$$S(P, Q) = \begin{pmatrix} 1 & 1 \\ -a & -b \end{pmatrix}$$

$$\begin{pmatrix} 1 & 1 \\ -a & -b \end{pmatrix} \begin{pmatrix} 1 & -1 \\ 0 & 1 \end{pmatrix}$$

Congruence number

$$P(X) = X - a, \quad Q(X) = X - b.$$

$$S(P, Q) = \begin{pmatrix} 1 & 1 \\ -a & -b \end{pmatrix}$$

$$\begin{pmatrix} 1 & 1 \\ -a & -b \end{pmatrix} \begin{pmatrix} 1 & -1 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ -a & a - b \end{pmatrix}.$$

Congruence number

$$P(X) = X - a, \quad Q(X) = X - b.$$

$$S(P, Q) = \begin{pmatrix} 1 & 1 \\ -a & -b \end{pmatrix}$$

$$\begin{pmatrix} 1 & 1 \\ -a & -b \end{pmatrix} \begin{pmatrix} 1 & -1 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ -a & a - b \end{pmatrix}.$$

\Rightarrow Congruence number $c(P, Q) = a - b$.

Congruence number

$$P(X) = X^2 + X + 1, \quad Q(X) = X - 1.$$

$$S(P, Q) = \begin{pmatrix} 1 & 1 & 0 \\ 1 & -1 & 1 \\ 1 & 0 & -1 \end{pmatrix}$$

Congruence number

$$P(X) = X^2 + X + 1, \quad Q(X) = X - 1.$$

$$S(P, Q) = \begin{pmatrix} 1 & 1 & 0 \\ 1 & -1 & 1 \\ 1 & 0 & -1 \end{pmatrix} \sim \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 2 & 2 & 3 \end{pmatrix}.$$

Congruence number

$$P(X) = X^2 + X + 1, \quad Q(X) = X - 1.$$

$$S(P, Q) = \begin{pmatrix} 1 & 1 & 0 \\ 1 & -1 & 1 \\ 1 & 0 & -1 \end{pmatrix} \sim \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 2 & 2 & 3 \end{pmatrix}.$$

⇒ Congruence number $c(P, Q) = 3$.

Congruence number

$$P(X) = X^2 + X + 1, \quad Q(X) = (X - 1)(X + 2) = X^2 + X - 2.$$

$$S(P, Q) = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & -2 & 1 \\ 0 & 1 & 0 & -2 \end{pmatrix}$$

Congruence number

$$P(X) = X^2 + X + 1, \quad Q(X) = (X - 1)(X + 2) = X^2 + X - 2.$$

$$S(P, Q) = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & -2 & 1 \\ 0 & 1 & 0 & -2 \end{pmatrix} \sim \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 3 & 0 \\ 2 & 1 & 0 & 3 \end{pmatrix}$$

Congruence number

$$P(X) = X^2 + X + 1, \quad Q(X) = (X - 1)(X + 2) = X^2 + X - 2.$$

$$S(P, Q) = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & -2 & 1 \\ 0 & 1 & 0 & -2 \end{pmatrix} \sim \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 3 & 0 \\ 2 & 1 & 0 & 3 \end{pmatrix}$$

\Rightarrow Congruence number $c(P, Q) = 3$.

(The resultant is 9.)

Next: Partial solution to our problem.

Congruence number

Theorem. Let $P, Q \in \mathbb{Z}[X]$.

Congruence number

Theorem. Let $P, Q \in \mathbb{Z}[X]$.

Let $r, s \in \mathbb{Z}[X]$ such that for the congruence number

$$\ell^n \mid c(P, Q) = rP + sQ.$$

Congruence number

Theorem. Let $P, Q \in \mathbb{Z}[X]$.

Let $r, s \in \mathbb{Z}[X]$ such that for the congruence number

$$\ell^n \mid\mid c(P, Q) = rP + sQ.$$

Suppose one of the following holds:

- Neither P nor Q has a multiple factor mod ℓ .
- P has no multiple factor mod ℓ and P and r are coprime mod ℓ .
- Q has no multiple factor mod ℓ and Q and s are coprime mod ℓ .

Congruence number

Theorem. Let $P, Q \in \mathbb{Z}[X]$.

Let $r, s \in \mathbb{Z}[X]$ such that for the congruence number

$$\ell^n \mid\mid c(P, Q) = rP + sQ.$$

Suppose one of the following holds:

- Neither P nor Q has a multiple factor mod ℓ .
- P has no multiple factor mod ℓ and P and r are coprime mod ℓ .
- Q has no multiple factor mod ℓ and Q and s are coprime mod ℓ .

Then there are $\alpha, \beta \in \overline{\mathbb{Z}}$ such that

- (i) $P(\alpha) = Q(\beta) = 0$ and
- (ii) $\alpha \equiv \beta \pmod{\ell^n}$.

Next: Complete solution to our problem.

Newton polygon method

Consider factorisations (over a big enough extension):

$$P(X) = \prod_{i=1}^u (X - \alpha_i) \text{ and } Q(X) = \prod_{j=1}^v (X - \beta_j).$$

Newton polygon method

Consider factorisations (over a big enough extension):

$$P(X) = \prod_{i=1}^u (X - \alpha_i) \text{ and } Q(X) = \prod_{j=1}^v (X - \beta_j).$$

Now take $Q(X + Y) = \prod_{j=1}^v (X - (\beta_j - Y))$, considered as a polynomial in X with coefficients in $\mathbb{Z}[Y]$.

Newton polygon method

Consider factorisations (over a big enough extension):

$$P(X) = \prod_{i=1}^u (X - \alpha_i) \text{ and } Q(X) = \prod_{j=1}^v (X - \beta_j).$$

Now take $Q(X + Y) = \prod_{j=1}^v (X - (\beta_j - Y))$, considered as a polynomial in X with coefficients in $\mathbb{Z}[Y]$.

Let $F(Y)$ be the resultant of $P(X)$ and $Q(X + Y)$ wrt. X .

$$\Rightarrow F(Y) = \pm \prod_{i=1}^u \prod_{j=1}^v (Y - (\beta_j - \alpha_i)).$$

Newton polygon method

Consider factorisations (over a big enough extension):

$$P(X) = \prod_{i=1}^u (X - \alpha_i) \text{ and } Q(X) = \prod_{j=1}^v (X - \beta_j).$$

Now take $Q(X + Y) = \prod_{j=1}^v (X - (\beta_j - Y))$, considered as a polynomial in X with coefficients in $\mathbb{Z}[Y]$.

Let $F(Y)$ be the resultant of $P(X)$ and $Q(X + Y)$ wrt. X .

$$\Rightarrow F(Y) = \pm \prod_{i=1}^u \prod_{j=1}^v (Y - (\beta_j - \alpha_i)).$$

\Rightarrow The roots of $F(Y)$ are the numbers we are looking for.

Newton polygon method

$$F(Y) = \pm \prod_{i=1}^u \prod_{j=1}^v (Y - (\beta_j - \alpha_i))$$

\Rightarrow The slopes of the Newton Polygon of $F(Y) \in \mathbb{Z}_\ell[Y]$ are the $v_\ell(\beta_j - \alpha_i)$ (normalisation: $v_\ell(\ell) = 1$).

Newton polygon method

$$F(Y) = \pm \prod_{i=1}^u \prod_{j=1}^v (Y - (\beta_j - \alpha_i))$$

\Rightarrow The slopes of the Newton Polygon of $F(Y) \in \mathbb{Z}_\ell[Y]$ are the $v_\ell(\beta_j - \alpha_i)$ (normalisation: $v_\ell(\ell) = 1$).

Lemma. $m := \lceil v_\ell(\beta_j - \alpha_i) \rceil$.

Then $\beta_j \equiv \alpha_i \pmod{\ell^n}$ if and only if $m \geq n$.

Newton polygon method

$$F(Y) = \pm \prod_{i=1}^u \prod_{j=1}^v (Y - (\beta_j - \alpha_i))$$

\Rightarrow The slopes of the Newton Polygon of $F(Y) \in \mathbb{Z}_\ell[Y]$ are the $v_\ell(\beta_j - \alpha_i)$ (normalisation: $v_\ell(\ell) = 1$).

Lemma. $m := \lceil v_\ell(\beta_j - \alpha_i) \rceil$.

Then $\beta_j \equiv \alpha_i \pmod{\ell^n}$ if and only if $m \geq n$.

Proof.

$$\begin{aligned} \beta_j \equiv \alpha_i \pmod{\ell^n} &\Leftrightarrow \beta_j - \alpha_i \in (\pi_K)^{\gamma_{K/\mathbb{Q}_\ell}(n)} \\ &\Leftrightarrow ev(\beta_j - \alpha_i) \geq \gamma_{K/\mathbb{Q}_\ell}(n) = e(n-1) + 1 \\ &\Leftrightarrow v(\beta_j - \alpha_i) \geq (n-1) + \frac{1}{e} \Leftrightarrow \lceil v(\beta_j - \alpha_i) \rceil \geq n. \quad \square \end{aligned}$$

Newton polygon method

$$F(Y) = \pm \prod_{i=1}^u \prod_{j=1}^v (Y - (\beta_j - \alpha_i))$$

\Rightarrow The slopes of the Newton Polygon of $F(Y) \in \mathbb{Z}_\ell[Y]$ are the $v_\ell(\beta_j - \alpha_i)$ (normalisation: $v_\ell(\ell) = 1$).

Lemma. $m := \lceil v_\ell(\beta_j - \alpha_i) \rceil$.

Then $\beta_j \equiv \alpha_i \pmod{\ell^n}$ if and only if $m \geq n$.

Proof.

$$\begin{aligned} \beta_j \equiv \alpha_i \pmod{\ell^n} &\Leftrightarrow \beta_j - \alpha_i \in (\pi_K)^{\gamma_{K/\mathbb{Q}_\ell}(n)} \\ &\Leftrightarrow ev(\beta_j - \alpha_i) \geq \gamma_{K/\mathbb{Q}_\ell}(n) = e(n-1) + 1 \\ &\Leftrightarrow v(\beta_j - \alpha_i) \geq (n-1) + \frac{1}{e} \Leftrightarrow \lceil v(\beta_j - \alpha_i) \rceil \geq n. \quad \square \end{aligned}$$

The biggest slope is the number solving the problem.

Conclusion

Let $P, Q \in \mathbb{Z}[X]$ be monic coprime polynomials.

The maximum n such that there are $\alpha, \beta \in \overline{\mathbb{Z}}$ satisfying

- (i) $P(\alpha) = Q(\beta) = 0$ and
- (ii) $\alpha \equiv \beta \pmod{\ell^n}$

can be computed

- in most cases by the congruence number
- always with the Newton polygon method.

Next: Modular Forms.

Computing modular forms

Let f be a newform (level N , weight k) with Fourier expansion:

$$f = f(z) = \sum_{m=1}^{\infty} a_m(f)q^m \text{ with } q = q(z) = e^{2\pi iz}.$$

Fact: All the $a_m(f)$ are integers of some number field.

Computing modular forms

Let f be a newform (level N , weight k) with Fourier expansion:

$$f = f(z) = \sum_{m=1}^{\infty} a_m(f)q^m \text{ with } q = q(z) = e^{2\pi iz}.$$

Fact: All the $a_m(f)$ are integers of some number field.

$\text{Gal}(\overline{\mathbb{Q}}/\mathbb{Q})$ naturally acts on the Fourier expansion.

$$\rightsquigarrow [f] := \mathbb{Z}\text{-span of } \text{Gal}(\overline{\mathbb{Q}}/\mathbb{Q}).f.$$

Computing modular forms

Let f be a newform (level N , weight k) with Fourier expansion:

$$f = f(z) = \sum_{m=1}^{\infty} a_m(f)q^m \text{ with } q = q(z) = e^{2\pi iz}.$$

Fact: All the $a_m(f)$ are integers of some number field.

$\text{Gal}(\overline{\mathbb{Q}}/\mathbb{Q})$ naturally acts on the Fourier expansion.

$$\rightsquigarrow [f] := \mathbb{Z}\text{-span of } \text{Gal}(\overline{\mathbb{Q}}/\mathbb{Q}).f.$$

Fact that makes computations possible:

$a_p(f)$ is a zero of the characteristic polynomial $P_{f,p} \in \mathbb{Z}[X]$ of the Hecke operator T_p acting on $[f]$.

$P_{f,p}$ is easy to compute!

Congruences of modular forms mod ℓ^n

$f = \sum_{m=1}^{\infty} a_m(f)q^m$ a newform (level N_f , weight k).

$g = \sum_{m=1}^{\infty} a_m(g)q^m$ a newform (level N_g , weight k).

Definition. f and g are congruent modulo ℓ^n if

$a_p(f) \equiv a_p(g) \pmod{\ell^n}$ for (almost) all primes p .

Congruences of modular forms mod ℓ^n

$f = \sum_{m=1}^{\infty} a_m(f)q^m$ a newform (level N_f , weight k).

$g = \sum_{m=1}^{\infty} a_m(g)q^m$ a newform (level N_g , weight k).

Definition. f and g are congruent modulo ℓ^n if

$a_p(f) \equiv a_p(g) \pmod{\ell^n}$ for (almost) all primes p .

If f and g are congruent mod ℓ^n , then

$P_{f,p}$ and $P_{g,p}$ have zeros which are congruent mod ℓ^n .

(Recall: $P_{f,p}$, $P_{g,p}$ characteristic polynomials of T_p on $[f]$ and $[g]$.)

Congruences of modular forms mod ℓ^n

$f = \sum_{m=1}^{\infty} a_m(f)q^m$ a newform (level N_f , weight k).

$g = \sum_{m=1}^{\infty} a_m(g)q^m$ a newform (level N_g , weight k).

Definition. f and g are congruent modulo ℓ^n if

$a_p(f) \equiv a_p(g) \pmod{\ell^n}$ for (almost) all primes p .

If f and g are congruent mod ℓ^n , then

$P_{f,p}$ and $P_{g,p}$ have zeros which are congruent mod ℓ^n .

(Recall: $P_{f,p}$, $P_{g,p}$ characteristic polynomials of T_p on $[f]$ and $[g]$.)

Some propositions (+ a very believable hypothesis)

\Rightarrow converse is true if compute 'enough' p (Sturm bound).

\rightsquigarrow Use congruence numbers/Newton polygon method!

Congruences of modular forms mod ℓ^n

Algorithm:

Congruences of modular forms mod ℓ^n

Algorithm:

$$c_2 := c(P_{f,2}, P_{g,2})$$

Congruences of modular forms mod ℓ^n

Algorithm:

$$c_2 := c(P_{f,2}, P_{g,2})$$

$$c_3 := c(P_{f,3}, P_{g,3})$$

Congruences of modular forms mod ℓ^n

Algorithm:

$$c_2 := c(P_{f,2}, P_{g,2})$$

$$c_3 := c(P_{f,3}, P_{g,3})$$

$$c_5 := c(P_{f,5}, P_{g,5})$$

... ■

Congruences of modular forms mod ℓ^n

Algorithm:

$$c_2 := c(P_{f,2}, P_{g,2})$$

$$c_3 := c(P_{f,3}, P_{g,3})$$

$$c_5 := c(P_{f,5}, P_{g,5})$$

...

⇒ **Upper bound** $u := \gcd(c_2 \cdot 2^\infty, c_3 \cdot 3^\infty, c_5 \cdot 5^\infty, \dots)$.

Prop. f and g are incongruent mod ℓ^m whenever $\ell^m \nmid u$.

Congruences of modular forms mod ℓ^n

Algorithm:

$$c_2 := c(P_{f,2}, P_{g,2})$$

$$c_3 := c(P_{f,3}, P_{g,3})$$

$$c_5 := c(P_{f,5}, P_{g,5})$$

...

⇒ **Upper bound** $u := \gcd(c_2 \cdot 2^\infty, c_3 \cdot 3^\infty, c_5 \cdot 5^\infty, \dots)$.

Prop. f and g are incongruent mod ℓ^m whenever $\ell^m \nmid u$.

Conversely, from congruence number/Newton method and Sturm bound get like this the maximum m such that (under hypothesis):

$$f \equiv g \pmod{\ell^n}.$$

Special case: Level N and Np

Weight 2, level $\Gamma_0(N)$, prime $p \nmid N$. Degeneracy maps:

$$S_2(N) \begin{array}{c} f(q) \mapsto f(q) \\ \xrightarrow{\hspace{2cm}} \\ f(q) \mapsto f(q^p) \end{array} S_2(Np).$$

Special case: Level N and Np

Weight 2, level $\Gamma_0(N)$, prime $p \nmid N$. Degeneracy maps:

$$S_2(N) \begin{array}{c} f(q) \mapsto f(q) \\ \xrightarrow{\hspace{2cm}} \\ f(q) \mapsto f(q^p) \end{array} S_2(Np).$$

Span: p -oldspace. Hecke operator \tilde{T}_p on oldspace

$$\begin{pmatrix} T_p & 1_d \\ -p \cdot 1_d & 0_d \end{pmatrix} \text{ with } T_p \text{ Hecke operator on } [f].$$

Let $\tilde{P}_{f,p}$ be the charpoly of \tilde{T}_p .

Special case: Level N and Np

Weight 2, level $\Gamma_0(N)$, prime $p \nmid N$. Degeneracy maps:

$$S_2(N) \begin{array}{c} f(q) \mapsto f(q) \\ \xrightarrow{\hspace{2cm}} \\ f(q) \mapsto f(q^p) \end{array} S_2(Np).$$

Span: p -oldspace. Hecke operator \tilde{T}_p on oldspace

$$\begin{pmatrix} T_p & 1_d \\ -p \cdot 1_d & 0_d \end{pmatrix} \text{ with } T_p \text{ Hecke operator on } [f].$$

Let $\tilde{P}_{f,p}$ be the charpoly of \tilde{T}_p .

Modify the algorithm at p :

Compute $c(\tilde{P}_{f,p}, P_{g,p})$ (instead of $c(P_{f,p}, P_{g,p})$).

Examples

Extract from Xavier's table:

N_1	i_1	N_2	i_2	lower bound	upper bound
71	2	71	1	$2 \cdot 3^2$	$2 \cdot 3^2$
109	3	109	1	2^2	2^2
155	4	155	2	2^4	2^4
233	3	233	1	3^3	3^3
613	2	613	1	$7 \cdot 47^2$	$7 \cdot 47^2$
785	2	785	1	7^3	7^3
1073	6	1073	3	$2 \cdot 17^2$	$2 \cdot 17^2$
1481	3	1481	1	$5^2 \cdot 2833$	$5^2 \cdot 2833$
1939	4	1939	3	$37^2 \cdot 4423$	$37^2 \cdot 4423$

Weak \neq Strong

Weight 2, level $\Gamma_0(71)$. Two classes of newforms both with coefficient field K : $[K : \mathbb{Q}] = 3$ and discriminant $d_K = 257$:

$$[f] : K \xrightarrow{\cong} \overline{\mathbb{Q}} \xrightarrow{\text{fixed}} \overline{\mathbb{Q}}_3 \quad \text{and} \quad [g] : K \xrightarrow{\cong} \overline{\mathbb{Q}} \xrightarrow{\text{fixed}} \overline{\mathbb{Q}}_3.$$

Weak \neq Strong

Weight 2, level $\Gamma_0(71)$. Two classes of newforms both with coefficient field K : $[K : \mathbb{Q}] = 3$ and discriminant $d_K = 257$:

$$[f] : K \xrightarrow{\cong} \overline{\mathbb{Q}} \xrightarrow{\text{fixed}} \overline{\mathbb{Q}}_3 \quad \text{and} \quad [g] : K \xrightarrow{\cong} \overline{\mathbb{Q}} \xrightarrow{\text{fixed}} \overline{\mathbb{Q}}_3.$$

Modulo 3, both f_1 and g_1 land in \mathbb{F}_3 , others in \mathbb{F}_9 .

$$f_1 \equiv g_1 \pmod{9} \text{ and } f_1 \not\equiv g_1 \pmod{27}$$

\Rightarrow Only one strong Hecke eigenform modulo 9.

Weak \neq Strong

Weight 2, level $\Gamma_0(71)$. Two classes of newforms both with coefficient field K : $[K : \mathbb{Q}] = 3$ and discriminant $d_K = 257$:

$$[f] : K \xrightarrow{\cong} \overline{\mathbb{Q}} \xrightarrow{\text{fixed}} \overline{\mathbb{Q}}_3 \quad \text{and} \quad [g] : K \xrightarrow{\cong} \overline{\mathbb{Q}} \xrightarrow{\text{fixed}} \overline{\mathbb{Q}}_3.$$

Modulo 3, both f_1 and g_1 land in \mathbb{F}_3 , others in \mathbb{F}_9 .

$$f_1 \equiv g_1 \pmod{9} \text{ and } f_1 \not\equiv g_1 \pmod{27}$$

\Rightarrow Only one strong Hecke eigenform modulo 9.

$\hat{\mathbb{T}}$:= Hecke algebra completed at 3.

The local factor $\hat{\mathbb{T}}_{\mathfrak{m}}$ belonging to f_1 and g_1 satisfies:

$$\hat{\mathbb{T}}_{\mathfrak{m}} \twoheadrightarrow \hat{\mathbb{T}}_{\mathfrak{m}} \otimes_{\mathbb{Z}_3} \mathbb{Z}/9\mathbb{Z} \cong \mathbb{Z}/9\mathbb{Z}[X]/(X^2) \xrightarrow{X \mapsto 0 \text{ or } X \mapsto 3} \mathbb{Z}/9\mathbb{Z}.$$

Hence: two weak Hecke eigenforms!

Level raising mod ℓ^n

Question. Given: f in level N , weight k , a prime $p \nmid N$ s.t.
 $\ell^n \mid c(P_{f,p}, X - (p + 1))$ or $\ell^n \mid c(P_{f,p}, X + (p + 1))$.

Is there g in level Np , weight k such that $f \equiv g \pmod{\ell^n}$?

Level raising mod ℓ^n

Question. Given: f in level N , weight k , a prime $p \nmid N$ s.t.
 $\ell^n \mid c(P_{f,p}, X - (p + 1))$ or $\ell^n \mid c(P_{f,p}, X + (p + 1))$.

Is there g in level Np , weight k such that $f \equiv g \pmod{\ell^n}$?

(Famous theorem by Ribet (Diamond, Taylor) for $n = 1$.)

Level raising mod ℓ^n

Question. Given: f in level N , weight k , a prime $p \nmid N$ s.t.
 $\ell^n \mid c(P_{f,p}, X - (p + 1))$ or $\ell^n \mid c(P_{f,p}, X + (p + 1))$.

Is there g in level Np , weight k such that $f \equiv g \pmod{\ell^n}$?

(Famous theorem by Ribet (Diamond, Taylor) for $n = 1$.)

Example. f in level 17, weight 2. Coefficients in \mathbb{Z} .

Level raising mod ℓ^n

Question. Given: f in level N , weight k , a prime $p \nmid N$ s.t.
 $\ell^n \mid c(P_{f,p}, X - (p + 1))$ or $\ell^n \mid c(P_{f,p}, X + (p + 1))$.

Is there g in level Np , weight k such that $f \equiv g \pmod{\ell^n}$?

(Famous theorem by Ribet (Diamond, Taylor) for $n = 1$.)

Example. f in level 17, weight 2. Coefficients in \mathbb{Z} .

$a_{59}(f) = 12$: congruence numbers

$$9 \parallel c(X - 12, X + (59 + 1)) = -72,$$

$$3 \parallel c(X - 12, X - (59 + 1)) = 48.$$

Level raising mod ℓ^n

Question. Given: f in level N , weight k , a prime $p \nmid N$ s.t.
 $\ell^n \mid c(P_{f,p}, X - (p + 1))$ or $\ell^n \mid c(P_{f,p}, X + (p + 1))$.

Is there g in level Np , weight k such that $f \equiv g \pmod{\ell^n}$?

(Famous theorem by Ribet (Diamond, Taylor) for $n = 1$.)

Example. f in level 17, weight 2. Coefficients in \mathbb{Z} .

$a_{59}(f) = 12$: congruence numbers

$$9 \parallel c(X - 12, X + (59 + 1)) = -72,$$
$$3 \parallel c(X - 12, X - (59 + 1)) = 48.$$

In level $17 \cdot 59$, weight 2, \exists 3 newforms g_1, g_2, g_3 s.t.

$$g_i \equiv f \pmod{3} \text{ for all } i = 1, 2, 3,$$

Level raising mod ℓ^n

Question. Given: f in level N , weight k , a prime $p \nmid N$ s.t.
 $\ell^n \mid c(P_{f,p}, X - (p + 1))$ or $\ell^n \mid c(P_{f,p}, X + (p + 1))$.

Is there g in level Np , weight k such that $f \equiv g \pmod{\ell^n}$?

(Famous theorem by Ribet (Diamond, Taylor) for $n = 1$.)

Example. f in level 17, weight 2. Coefficients in \mathbb{Z} .

$a_{59}(f) = 12$: congruence numbers

$$9 \parallel c(X - 12, X + (59 + 1)) = -72,$$
$$3 \parallel c(X - 12, X - (59 + 1)) = 48.$$

In level $17 \cdot 59$, weight 2, \exists 3 newforms g_1, g_2, g_3 s.t.

$$g_i \equiv f \pmod{3} \text{ for all } i = 1, 2, 3,$$

but there is no i s.t. $g_i \equiv f \pmod{9}$!

Level raising mod ℓ^n

Does a weaker statement hold?

Level raising mod ℓ^n

Does a weaker statement hold?

Let g_1, \dots, g_r be all newforms in $S_2(\Gamma_0(Np))$.

Let ℓ^{n_i} be the highest power of ℓ such that

$$g_i \equiv f \pmod{\ell^{n_i}} \text{ for } i = 1, \dots, r.$$

Put $n := n_1 + \dots + n_r$.

Put $c := c(P_{f,p}, X^2 - (p+1)^2)$.

Level raising mod ℓ^n

Does a weaker statement hold?

Let g_1, \dots, g_r be all newforms in $S_2(\Gamma_0(Np))$.

Let ℓ^{n_i} be the highest power of ℓ such that

$$g_i \equiv f \pmod{\ell^{n_i}} \text{ for } i = 1, \dots, r.$$

Put $n := n_1 + \dots + n_r$.

Put $c := c(P_{f,p}, X^2 - (p+1)^2)$.

Question. Is n at least as big as (or even equal to) the ℓ -valuation of c ?

THE END